**COMP 3270 Homework 1**

100 points. Due by **11:59pm (midnight) on Wednesday, September 8th, 2021**

Instructions:

1. This is an individual assignment. You should do your own work. Any evidence of copying will result in a zero grade and additional penalties/actions.
2. Submissions not handed on the due date and time **will not** be accepted unless prior permission has been granted or there is a valid and verifiable excuse.
3. Think carefully; formulate your answers, and then write them out concisely using English, logic, mathematics and pseudocode (no programming language syntax).
4. Type your final answers in this Word document and submit online through Canvas.
5. Don’t turn in handwritten answers with scribbling, cross-outs, erasures, etc. If an answer is unreadable, it will earn zero points. **Neatly and cleanly handwritten submissions are also acceptable**.

1. (3 points) *Computational problem solving: Estimating problem solving time:* Suppose there are three algorithms to solve a problem- a O(n) algorithm (A1), a O(nlogn) algorithm (A2) and a O(n2) algorithm (A3) where log is to the base 2. Using the techniques and assumptions in slide set L2-Buffet(SelectionProblem).ppt, determine how long in seconds it will take for each algorithm to solve a problem of size 200 million. You must show your work to get credit, i.e., a correct answer without showing how it was arrived at will receive zero credit.

**Machine Clock Speed = 4 x 109 clock cycles per second OR 2 x 107 steps per second**

**200 million = 2 x 108**

**O(n) algorithm (A1) = 2 x 108 / 2 x 107 = 10 seconds**

**O(nlogn) algorithm (A2) = [ (2 x 108) log­­­­­­­2(2 x 108) ] / (2 x 107) = (5.5 x 109) / ( 2x 107) = 275.75 seconds**

**O(n2) algorithm (A3) = [ (2 x 108)2 ] / (2 x 107) = (4 x 1016) / (2 x 107) = 2 x 109 seconds**

2. (6 points) *Computational problem solving: Problem specification*

Suppose you are asked to develop a mobile application to provide turn by turn directions on a smartphone to an AU parking lot in which there are at least five empty parking spots nearest to a campus building that a user selects. Assume that you can use the Google Map API for two functions (only) ─ display campus map on the phone so user can select a campus building, and produce turn-by-turn directions from a source location to a destination location ─ where any location in the map is specified as a pair (latitude, longitude). Also assume that there is an application called AUparking that you can query to determine the # of vacant spots in any parking lot specified as a pair (latitude, longitude). Specify the problem to a level of detail that would allow you to develop solution strategies and corresponding algorithms: State the problem specification in terms of (1) inputs, (2) data representation and (3) desired outputs; no need to discuss solution strategies.

**1. Inputs: Selected nearest campus building destination**

**2. Data Representation:**

**- Vacant spots available**

**- Locations (user’s and destination)**

**3. Output: User will have access to turn-by-turn directions to a selected parking lot (that is considered “vacant”) based off a campus building.**

3. (5 points) *Computational problem solving: Developing strategies*

Explain a correct and efficient **strategy** to check what the maximum difference is between any pair of numbers in an array containing n numbers. Your description should be such that the strategy is clear, but at the same time the description should be at the level of a strategy, not an algorithm. Then state the total number of number pairs any algorithm using the strategy “compute the difference between every number pair in the array and select that pair with the largest difference” will have to consider as a function of n.

**Strategy: The first step in determining the max difference between any pairs in an array would first, sort the array (MergeSort could possibly be the fastest method since it computes at a O(nlogn) rate). After sorting, you would have to index the last element and the first. Then finally, you would subtract those two elements to find the maximum difference.**

4. (7 points) *Computational problem solving: Understanding an algorithm and its strategy by simulating it on a specific input:*

Understand the following algorithm. Simulate it mentally on the following four inputs, and state the outputs produced (value returned) in each case: (a) A: [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]; (b) A: [─1, ─2, ─3, ─4, ─5, ─6, ─7, ─8, ─9, ─10], ; (c) A: [0, 0, 0, 0, 0, 0, 0, 0, 0, 0], (d) A: [─1, 2, ─3, 4, ─5, 6, 7, ─8, 9, ─10].

Algorithm-1 (A:array[1..n] of integer)

sum, max: integer

1 sum = 0

2 max = 0

3 for i = 1 to n

4 sum = 0

5 for j = i to n

6 sum = sum + A[j]

7 if sum > max then

8 max = sum

9 return max

Output when input is array (a) above:

**55**

Output when input is array (b) above:

**0**

Output when input is array (c) above:

**0**

Output when input is array (d) above:

**11**

What does the algorithm return when the input array contains all negative integers?

**0. This is due to the sum not being greater than the max value.**

What does the algorithm return when the input array contains all non-negative integers?

**A positive integer of the sum**

What does the algorithm return when the input array contains negative, zero and positive integers?

**It will return the max sum of any of the numbers.**

5. (9 points) *Computational problem solving: Calculating approximate complexity:*

Using the approach described in class (L5-Complexity.pptx), calculate the approximate complexity of Algorithm-1 above by filling in the table below.

|  |  |
| --- | --- |
| Step | Big-Oh complexity |
| **1 sum = 0** | **O(1)** |
| **2 max = 0** | **O(1)** |
| **3 for i = 1 to n** | **O(n)** |
| **4 sum = 0** | **O(1)** |
| **5 for j = i to n** | **O(n2)** |
| **6 sum = sum + A[j]** | **O(n2)** |
| **7 if sum > max then** | **O(n2)** |
| **8 max = sum** | **O(n2)** |
| **9 return max** | **O(1)** |
| **Complexity of the algorithm** | **O(n2) since it is the biggest complexity in all of code** |

6. (18 points) Calculate the detailed complexity T(n) of Algorithm-1. Fill in the table below, then determine the expression for T(n) and simplify it to produce a polynomial in n.

|  |  |  |
| --- | --- | --- |
| Step | Cost of each execution | Total # of times executed |
| 1 | **1** | **1** |
| 2 | **1** | **1** |
| 3 | **1** | **(n+1)** |
| 4 | **1** | **n** |
| 5 | **1** |  |
| 6 | **6** | **N2** |
| 7 | **3** | **N2** |
| 8 | **3** | **N2** |
| 9 | **2** | **1** |

T(n) = **1 + 1 + (n + 1) + n + (n-i) + 6n2 + 3n2 + 3n2 + 2**

**= 13n2 + 3n + 4**

7. (5 points) *Computational problem solving: Proving correctness/incorrectness:*

Is the algorithm below correct or incorrect? Prove it! It is supposed to count the number of all identical integers that appear consecutively in a file of integers. E.g., if f contains 1 2 3 3 3 4 3 5 6 6 7 8 8 8 8 then the correct answer is 9

Count(f: input file)

count, i, j : integer //local variables

count=0

while end-of-file(f)=false

i=read-next-integer(f)

if end-of-file(f)=false then

j=read-next-integer(f)

if i=j then count=count+1

return count

**Counterexample:**

**Output = 6**

**Correct Output = 9**

**The above algorithm is incorrect due to the fact that a wrong comparison is made between the integers; the count only compares the next integer with the current integer and does not count the current integer. If it were to be correct, the algorithm must include how it includes the current integer in its count.**

8. (10 points) *Computational problem solving: Proving correctness:* Complete the proof by contradiction this algorithm to compute the Fibonacci numbers is correct.

function fib(n)

1. if n=0 then return(1)

2. if n=1 then return(1)

3. last=1

4. current=1

5. for i=2 to n do

6. temp=last+current

7. last=current

8. current=temp

9. return(current)

1. Assume the algorithm is incorrect.
2. Fibonacci numbers are defined as F0=1, F1=1, Fi=Fi-1+Fi-2 for i>1.
3. So the assumption in (1) implies that there is at least one input parameter n=k, k≥0, for which the algorithm will produce an incorrect answer.
4. **Base cases: if n = 0 or n = 1, it will return 1.**

So in both cases the algorithm returns the correct answer.

1. This implies that there has to be at least one integer k>1, so that when n=k the algorithm does not return the correct answer Fk=Fk-1+Fk-2.
2. When n=k and k>1 \_\_\_\_\_**steps 1 and 2 will not be executed because k>1**\_\_\_\_, and steps 3-9 will be executed.
3. If k=2, the for loop in steps 5-8 will be executed exactly once. By step 6, temp = last + current = 1 + 1 = F0 + F1. Then step 7 updates last to be equal to current = F1. Step 7 updates current to be equal to temp which is F0 + F1. So the value returned in step 9 is current = F0 + F1 = F2. This is the correct answer. So the k for which the algorithm fails must be greater than 2.
4. If k=3, \_\_\_\_\_**Steps 1 and 2 will not execute because k>1. The for loop itself will execute 2 times, with each time returning the correct current value.**\_\_\_\_\_\_\_\_\_\_\_
5. But if k= 4, \_\_\_\_\_\_\_\_\_**Steps 1 and 2 will still not execute, but this time the for loop will execute 3 times, with each time still returning the correct current value.**\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_
6. The above argument can be repeated to show that \_\_\_\_\_\_\_\_**Steps 1 and 2 will not be executed, thus will continue to the for loop, where it will execute k-1 times, and return the correct current values.** \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_.
7. That is, for all k > 1 the algorithm returns the correct k-th Fibonacci number.
8. So there is no k for which the algorithm will return a value not equal to Fk-1+Fk-2. This contradicts (3).
9. Therefore, the algorithm must be correct.

9. (a) (6 points) *Computational problem solving: Algorithm design:* Describe a recursive algorithm to reverse a string that uses the strategy of swapping the first and last characters and recursively reversing the rest of the string. Assume the string is passed to the algorithm as an array A of characters, A[p…q], where the array has starting index p and ending index q, and the length of the string is n=q–p+1. The algorithm should have only one base case, when it gets an empty string. Assume you have a swap(A[i],A[j]) function available that will swap the characters in cells i and j. Write the algorithm using pseudocode without any programming language specific syntax. Your algorithm should be correct as per the technical definition of correctness.

**Function reverseString (A:array[p…q] of type int)**

**If ( p < A.length / 2 || q > A.length /2 )**

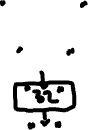
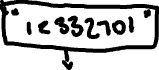
**A = swap ( A[p + 1], A[q - 1] )**

**Else**

**Return A**

**Return reverseString(A[ I + 1…j - 1])**

(b) (8 points) Draw your algorithm’s recursion tree on input string “i<33270!”- remember to show inputs and outputs of each recursive execution including the execution of any base cases.



10. (10 points) *Computational problem solving: Proving correctness:*

Function g (n: nonnegative integer)

if n ≤ 1 then return(n)

else return(5\*g(n-1) – 6\*g(n-2))

Prove by induction that algorithm g is correct, if it is intended to compute the function 3n-2n for all n ≥ 0.

Base Case Proof:

**N = 0 ----🡪 30 – 20 = 0.**

**N = 1 ---🡪 31 – 21 = 1.**

**From the above statements that are correct, it is noted that the base cases have been proven to be correct since when n = 0, it returns n (in this case 0) when n <= 1. And also when n = 1, it returns n (in this case 1) when n <= 1.**

Inductive Hypothesis:

**It is implied that for all n >= 0, that the output will return 3n-2n.**

Inductive Step:

**Because n = 0 and n = 1 have already been determined that it is correct, next we must determine the output for values greater than 1 and also n+1.**

11. (13 points) *Computational problem solving: Proving correctness:* The algorithm of Q.8 can also be proven correct using the Loop Invariant method. The proof will first show that it will correctly compute F0 & F1 by virtue of lines 1 and 2, and then show that it will correctly compute Fn, n>1, using the LI technique on the for loop. For this latter part of the correctness proof, complete the Loop Invariant below by filing in the blanks. Then complete the three parts of the rest of the proof.

Loop Invariant:

Before any execution of the for loop of line 5 in which the loop variable i=k, 2≤k≤n, the variable last will contain \_\_\_**fk-2**\_\_\_\_ and the variable current will contain \_\_**fk-1**\_\_\_\_\_\_\_.

Initialization:

**The loop initializes at i = 2, last = 1, current = 1. Loop invariant holds.**

Maintenance:

**This will require us to prove that the loop invariant holds until the ith execution. The loop invariant will hold at the end of the loop, where current will be Fk.**

Termination:

**It is determined that by finding out that the loop invariant holds after the last execution, it is therefore correct.**